

10/23/24

Lecture 14: Statistical Query Model (cont'd)

SQ dimension (recipe for supervised problems)

Def: Class of functions $\mathcal{F} = \{f: \mathbb{R} \rightarrow [-1, 1]\}$ has

SQ dimension $\geq D$ w.r.t η if $\exists f_1, \dots, f_D \in \mathcal{F}$
s.t. $\forall i \neq j$

$$\left| \mathbb{E}_{x \sim \eta} [f_i(x) f_j(x)] \right| \leq \frac{1}{D}.$$

Thm: If \mathcal{F} has SQ dimension D , then CSQ learning requires tolerance τ or $\mathcal{O}(D\tau^2)$ queries

Pf: For convenience, define $\langle f, g \rangle \triangleq \mathbb{E}[f(x)g(x)]$.

wLOG let $\mathcal{F} = \{f_1, \dots, f_D\}$.

Given query $\phi: \mathbb{R}^d \rightarrow [-1, 1]$, let

$$A^+ \triangleq \{f \in \mathcal{F} : \langle f, \phi \rangle \geq \tau\}$$

By Cauchy-Schwarz:

$$\begin{aligned} \left\langle \phi, \sum_{f \in A^+} f \right\rangle^2 &\leq \underbrace{\|\phi\|^2}_{\leq 1} \cdot \left\| \sum_{f \in A^+} f \right\|^2 \\ &\leq \sum_{f \in A^+} \|f\|^2 + \frac{|A^+|(|A^+|-1)}{D} \\ &\leq \frac{|A^+|^2}{D} + |A^+| \end{aligned}$$

But $\langle \phi, \sum_{f \in A^+} f \rangle \geq \tau |A^+|$ by defn., so

$$\tau^2 |A^+|^2 \leq \frac{|A^+|^2}{D} + |A^+|$$

$$\Rightarrow |A^+| \leq \frac{D}{D\tau^2 - 1} \leq O\left(\frac{1}{\tau^2}\right)$$

Similarly, for $A^- \triangleq \{f \in \mathcal{F} : \langle f, \phi \rangle \leq -\tau\}$,
 $|A^-| \leq O\left(\frac{1}{\tau^2}\right)$.

So regardless of ϕ , all but $O\left(\frac{1}{\tau^2}\right)$ many f 's consistent w/ the answer 0, so
need $\Omega(D\tau^2)$ queries. □

SQ dimension bound for 1-hidden-layer MCNs:

(Gel - Gollakota - Jin - Karmalkar - Klivans '20):

Let $S \subseteq [d]$ be of size $m = \lg_2 k$.

Given $w \in \{\pm 1\}^m$, define $w^{(S)} \in \mathbb{S}^{d-1}$ by

$$w_i^{(S)} = \begin{cases} w_i / \sqrt{m} & \text{if } i \in S \\ 0 & \text{o.w.} \end{cases}$$

Define $f_S : \mathbb{R}^d \rightarrow \mathbb{R}$ by

$$f_S(x) \triangleq \sum_{w \in \{\pm 1\}^m} \left(\prod_{i=1}^m w_i \right) \cdot \text{relu}(\langle w^{(S)}, x \rangle)$$

(Note: this is one-hidden-layer MLP w/ size $2^m = k$)

Claim: If γ is sign-symmetric dist. over \mathbb{R}

(i.e. for any x , x and $-x$ are equally likely under γ), then SQ dim of $\{f_S\}$ is $\geq \binom{d}{m}$.

Pf: For any distinct S, T of size m ,

$$\text{w.t.s } \langle f_S, f_T \rangle_\gamma = 0$$

For any $z \in \{\pm 1\}^n$, note that

$$f_S(x \odot z) = \sum_w \prod_i w_i \cdot \underbrace{\text{relu}(\langle w^{(S)}, x \odot z \rangle)}_{\text{relu}(\langle w^{(S)} \odot z, x \rangle)}$$

$$= \sum_w \prod_i w_i \cdot z_S \cdot \text{relu}(\langle w^{(S)}, x \rangle)$$

$$= z_S \cdot f_S(x).$$

by sign symmetry

$$\begin{aligned} \mathbb{E}_{x,z} \left[f_S(x \odot z) f_T(x \odot z) \right] &= \mathbb{E}_x \mathbb{E}_z \left[f_S(x) f_T(x) z_S z_T \right] \\ &= \mathbb{E}_x \left[f_S(x) f_T(x) \right] \cdot \underbrace{\mathbb{E}_z [z_S z_T]}_{\text{if } S \neq T}. \quad \square \end{aligned}$$

NB: Technically, also have to make sure $\{f_S\}$ are nonzero functions! See paper for this calculation (Hermite analysis).

(Diakonikolas - Kane - Kontonis - Zarifis '20)

Showed stronger lbd ($d^{-2(k)}$ queries or $\frac{2^{-\text{poly}(n)}}{\text{tolerance}}$)

using different instance of approximately orthogonal functions: given 2D subspace $U \subseteq \mathbb{R}^d$,

$$f_U(x) = h(\langle u_1, x \rangle, \langle u_2, x \rangle)$$

where $h(a, b) = \sum_{i=1}^k (-1)^i \langle \left(\cos\left(\frac{2\pi}{k}\right), \sin\left(\frac{2\pi}{k}\right) \right), (a, b) \rangle$.

and u_1, u_2 are basis for U .

idea: $\deg \leq \frac{k+1}{2}$ Hermite coeffs of f_u are 0, so

$$\langle f_U, f_V \rangle_{N(0, \text{Id})} = \langle f_U^{>\frac{k+1}{2}}, f_V^{>\frac{k+1}{2}} \rangle$$

$$\leq \left(\text{"Correlation}(U, V)" \right)^{(k)}.$$

So if we take a packing of the space of 2^d subspaces, we get the desired bound on the SQ dimension.

Statistical dimension :

Def: Set of distributions T has statistical dimension $\geq \Delta$ w.r.t. reference distribution D and avg. correlation γ if $H(T)$ of size $\geq |T^*|/\Delta$,

$$P_D(T) \triangleq \frac{1}{|T|^2} \sum_{D_i, D_j \in T} \langle D_i, D_j \rangle_D \leq \gamma,$$

where $\langle D_i, D_j \rangle_D = \bigoplus_{x \sim D} \left[\left(\frac{D_i(x)}{D(x)} - 1 \right) \left(\frac{D_j(x)}{D(x)} - 1 \right) \right]$

Thm: If T^* has statistical dimension $\geq \Delta$, then learning dist's in T^* w/ SQ alg. requires either $\Omega(\Delta)$ queries or $\sqrt{\gamma}$ tolerance.

PF: Given query $\phi: \mathbb{R}^d \rightarrow \mathbb{R}$, define

$$A^+ \stackrel{\Delta}{=} \left\{ D' \in T^* \text{ s.t. } \underset{D'}{\mathbb{E}}[\phi] \geq \underset{D}{\mathbb{E}}[\phi] + \sqrt{\gamma} \right\}$$

$$A^- \stackrel{\Delta}{=} \left\{ D' \in T^* \text{ s.t. } \underset{D'}{\mathbb{E}}[\phi] \leq \underset{D}{\mathbb{E}}[\phi] - \sqrt{\gamma} \right\},$$

i.e. set of dist's in T^* s.t. answering ϕ w/

$\underset{D}{\mathbb{E}}[\phi]$ is incorrect. w.t.s $|A^+|, |A^-|$ small.

$$\begin{aligned} \underset{x \sim D}{\mathbb{E}} \left[\phi(x) \cdot \sum_{D' \in A^+} \left\{ \frac{D'(x)}{D(x)} - 1 \right\} \right]^2 &\leq \underbrace{\underset{x}{\mathbb{E}} \left[\phi(x)^2 \right]}_{\leq 1} \cdot \underbrace{\mathbb{E}_x \left[\left(\sum_{D' \in A^+} \frac{D'(x)}{D(x)} - 1 \right)^2 \right]}_{= \sum_{D', D'' \in A^+} \langle D', D'' \rangle} \\ &\leq |A^+|^2 \cdot P_D(A^+) \end{aligned}$$

Remains to lower bound

$$\underset{x \sim D}{\mathbb{E}} \left[\phi(x) \cdot \sum_{D' \in A^+} \left\{ \frac{D'(x)}{D(x)} - 1 \right\} \right]^2. \quad (\dagger)$$

$$\text{Note: } \underset{x \sim D}{\mathbb{E}} \left[\phi(x) \cdot \left\{ \frac{D'(x)}{D(x)} - 1 \right\} \right] = \underset{x \sim D}{\mathbb{E}} [\phi(x)] - \underset{x \sim D}{\mathbb{E}} [\phi(x)]$$

$$\geq \sqrt{8},$$

so $(\delta) \geq |A^+|^2 \cdot \gamma$, so

$$f_\delta(A^+) \geq \gamma,$$

implying $|A^+| / |T^*| \leq 1/\Delta$.

Similarly, $|A^-| / |T^*| \leq 1/\Delta$.

So need $\mathcal{O}(\Delta)$ queries b/c each one only rules out a $1/\Delta$ fraction. \square

(Bonus material)

Moment-matching and SQ:

Consider P_V and $P_{V'}$ (defined in slides)

for moment-matching distribution A . WLOG

Suppose $V = (1, 0, \dots, 0)$, $V' = (\cos \theta, \sin \theta, 0, \dots, 0)$.

Claim : Let \mathcal{Q} be dist over \mathbb{R}^d

given by taking $X \sim P_V$ and outputting $V' \cdot X$.

$$\text{Then } \langle Q, Q \rangle_{N(0,1)} \leq (\cos \theta)^{2(n+1)} \langle A, A \rangle_{N(0,1)}$$

Pf: $Q(x) = \int P_v(\vec{x}) d\vec{x}$

$$\vec{x} \cdot (\vec{x}, v') = x$$

$$(x_1, z)$$

$$\langle (x_1, z), v' \rangle = x$$

$$\Downarrow$$

$$x_1 \cos \theta + z \sin \theta = x$$

$$x = \frac{x_1 + z}{\sqrt{2}}$$

Gaussian

Scaled convolution of A and Gaussian

$$= \int A(x_1) G(z_1) dx_1 dz_1,$$

$$(x_1, z_1):$$

$$x_1 \cos \theta + z \sin \theta = x$$

$$\int A(x) \gamma(x) \underset{\Theta}{=} \bigcup_{\Theta} A(x)$$

"Ornstein-Uhlenbeck noise operator"

(note: to sample from $U_\theta A$, draw $x \sim A$ and output $\cos \theta \cdot x + \sin \theta g$ for $g \sim N(0, 1)$)

Write $\frac{A(x)}{G(x)} = \sum_{i=0}^{\infty} a_i \underbrace{H e_i(x)}_{\text{probabilist's Hermite polynomial}} \cdot \frac{1}{\sqrt{i!}}$

Note: $a_i = \mathbb{E}_{x \sim N(0, 1)} \left[\frac{A(x)}{G(x)} H e_i(x) \cdot \frac{1}{\sqrt{i!}} \right]$

$\stackrel{b/c \text{ moment matching}}{=} \int A(x) H e_i(x) \cdot \frac{1}{\sqrt{i!}} dx$ orthogonality

$\stackrel{\downarrow}{=} \int G(x) H e_i(x) \cdot \frac{1}{\sqrt{i!}} dx = 0$

Mehler's lemma: $U_\theta (H e_i \cdot G)(x) =$

$$\cos^i(\theta) H e_i(x) G(x),$$

i.e. (Hermite \times Gaussian density) is eigenfunction of noise operator

So $U_\theta A(x) = \sum_{i=0}^{\infty} a_i \cos^i(\theta) H e_i(x) G(x) \frac{1}{\sqrt{i!}}$

$\stackrel{(\text{by } (1))}{=} \sum_{i=m+1}^{\infty} a_i \cos^i(\theta) H e_i(x) G(x) \frac{1}{\sqrt{i!}}$

$$\langle Q, Q \rangle_{N(0,1)} = \langle V_\theta A, V_\theta A \rangle_{N(0,1)}$$

$$= \sum_{i=m+1}^{\infty} a_i^2 \cos^{2i}(\theta)$$

$$\leq \cos^{2m+2}(\theta) \cdot \underbrace{\sum_{i=m+1}^{\infty} a_i^2}_{\text{``"}}$$

$$\langle A, A \rangle_{N(0,1)}. \square$$

We conclude that

$$\langle P_V, P_V \rangle_{N(0, \text{Id})} = \langle A, V_\theta A \rangle_{N(0,1)}$$

by the above

$$\leq \langle A, A \rangle_{N(0,1)}^{1/2} \cdot \langle V_\theta A, V_\theta A \rangle_{N(0,1)}^{1/2}$$

$$\leq \underbrace{\cos^{m+1}(\theta)}_{\text{exp. decaying in } m} \langle A, A \rangle_{N(0,1)}. \quad \square$$