

Lecture 20: Statistical Physics I: Intro to graphical models and variational inference

1 Wrapping up Cryptographic Hardness

1.1 Modern Results

Lemma 1 ([DV21]). *Assuming security of Goldreich's pseudorandom generator, MLP's of depth three are hard to learn under Gaussian inputs.*

The lifting method used in the proof of this result (see Lec. 19) can be extended to give SQ and membership query lower bounds for real valued continuous functions by exhibiting functions close to being boolean except for a certain pathological space of inputs [CGKM22]. Also, by slightly increasing the neural network complexity, hardness for depth 4 MLP's persists in the smooth analysis setting which is not fully adversarial (i.e. start with worst-case input and perturb parameters with Gaussian noise) [DSV23].

Lemma 2 ([BRST21]). *Under the assumption that there are no polynomial time quantum algorithms for certain (worst-case) lattice problems, learning mixtures of Gaussians is hard.*

Proof. The assumption of the lemma is related to learning with errors/solving a noisy system of linear equations over finite fields. The mixtures of Gaussians problem is the continuous analog to learning with errors problem: you can think of the parallel pancakes problem (which showed learning Gaussian mixtures was hard in SQ) as solving a noisy system of linear equations modulo 1. It turns out that learning with errors over finite fields is hard under the same assumption on quantum algorithms, so some kind of reduction as in [GVV22] from finite fields to the continuous analog can show hardness of mixtures of Gaussians. \square

1.2 Recap

- There are strong connections between cryptography/average-case complexity (learning parity or learning with errors) and machine learning theory, but because distributions in cryptography tend to be discrete, immediate reductions to ML theory tend to be pathological and unnatural or require agnostic noise to simulate.

- As a result, hardness results are restricted to models like statistical query learning. However, these are surprisingly robust: usually a lower bound for SQ is accompanied by a lower bound for cryptographic hardness a few years later. These lower bounds to restricted models are easier to prove than via reductions due to established recipes like the SQ dimension, moment matching, etc.
- Ultimately, these recipes can show that in the SQ model it is not just techniques like parity and low-degree approximation or Gaussian mixtures and moment methods/dimensionality reduction that don't work in the learning setting, but that any algorithm in the restricted setting fails to efficiently learn. In a sense, it is because these "signature methods" like moment matching or low-degree approximation are "optimal" in the SQ model, so if these methods fail, then no other method can work.
- Lower bounds help modulate the learning model: i.e., determining whether the model is too general for positive learnability results requires the lens of computational (time/compute complexity) lower bounds instead of statistical (data size) lower bounds. For example, distribution-free agnostic learning was once a target reasonable learnability setting, but it turned out to be too computationally hard.

2 Physics, inference, and sampling

Many problems in machine learning and statistics can be framed as the following inference setting. Suppose we have a prior distribution \mathcal{D} where we draw a signal $X \sim \mathcal{D}$. The signal undergoes some kind of noisy channel, and the observed "noisy measurement" is Y . The goal is to understand how the knowledge of Y induces a posterior belief on the original X . For example,

- **Learning neural networks/teacher-student setting:** If X denotes the weights of a neural network F sampled from some complicated prior, Y is the set $\{(x_i, y_i) : x_i \sim \mathcal{N}(0, \text{Id}_d), y_i = F(x_i)\}$ of input and output pairs of F , where the input is sampled from a multivariate Gaussian.
- **Denosing:** If X is an image from a distribution \mathcal{D} , Y is the image with added noise terms to the pixels. Can we recover the original X ? This setting turns out to be the core primitive of diffusion models for generative modeling.

From Bayes Rule, we can immediately express the posterior on X given Y :

$$P(X = x|Y = y) = \frac{P(Y = y|X = x)P(X = x)}{P(Y = y)}. \quad (1)$$

$P(X = x)$ is our prior on X , $P(Y = y|X = x)$ is the likelihood given a model of the noisy channel, and $P(Y = y)$ is the normalization constant, also known as the *partition function/evidence*. We can write this in a suggestive manner:

$$P(X = x|Y = y) = \frac{1}{Z} \cdot \exp(-\mathcal{E}(x)), \quad (2)$$

where the *energy function* $\mathcal{E}(x) = -\log(P(Y = y|X = x)P(X = x))$. This is highly reminiscent of the *Boltzmann energy distribution* from statistical physics. Finding the minimal energy corresponds to finding the signal X at which the posterior is maximized.

Further details on the background used in this and subsequent sections can be found in [MM09], [KZ22], [Mon11], [WJ08], and [Mon14].

3 Undirected graphical models with pairwise interactions

Consider the energy model given by

$$\mathcal{E}(x) = - \sum_{(i,j) \in \mathcal{F}} \log \psi_{ij}(x_i, x_j), \quad (3)$$

with \mathcal{F} some family of edges/subset of $[n] \times [n]$, inputs $x \in \{\pm 1\}^n$, and *compatibility functions* $\psi_{ij} : \{\pm 1\}^2 \rightarrow \mathbb{R}$. The posterior distribution is also known as the *Gibbs measure* μ over $\{\pm 1\}^n$, where $\mu = \frac{1}{Z} \cdot \exp(-\mathcal{E}(x))$ and the partition function $Z = \sum_{x \in \{\pm 1\}^n} \exp(-\mathcal{E}(x))$.

3.1 Ising Model

Suppose we have the *compatibility functions*

$$\psi_{ij} = \exp(-\beta x_i x_j A_{ij}) \quad (4)$$

for some matrix A such that $A_{ii} = 0$ (WLOG, since constants can be absorbed in the normalization) and A is symmetric. In this setting β can be interpreted as the

inverse temperature and A as the *Hamiltonian* or *interaction matrix* of the system. It follows from the definitions above that the Gibbs measure is given by

$$\mu \propto \exp\left(-\frac{\beta}{2}x^\top Ax\right). \quad (5)$$

For some limiting cases with the inverse temperature β , notice the high temperature limit $\beta \rightarrow 0$ corresponds to a uniform distribution over $\{\pm 1\}^n$, while the low temperature limit $\beta \rightarrow \infty$ corresponds to the uniform distribution over $x \in \{\pm 1\}^n$ such that $x^\top Ax$ is maximal.

3.2 Markov Property

We can also interpret A as the (negative) adjacency matrix of a weighted graph G . Given $i \in [n]$, define $\partial_i = \{j \in [n] : (i, j) \in \mathcal{F}\}$ for some edgeset \mathcal{F} corresponding to A ; i.e., the neighborhood of i where $A_{ij} \neq 0$.

The *Markov property* states that for $S \subseteq [n]$ such that $[n]/S$ consists of two disjoint pieces in G , then conditioning on some assignment of spins/ x values on S , the resulting marginals on the two disjoint pieces are independent. For example, if $S = \partial_i$, then G is broken into i and everything not connected to i . Then the marginal on i conditioned on some assignment to ∂_i is independent of the rest of the graph, i.e.,

$$P(x_i = \sigma | x_{\partial_i} = s) \propto \exp\left(-\beta \sum_{j \in \partial_i} A_{ij} s_j \sigma\right). \quad (6)$$

4 Variational Inference

In the types of inference problems detailed above, there are two fundamental tasks: computing the normalization constant/partition function Z , and sampling from the Gibbs measure μ . For discrete but exponentially sized domains, computing Z is incredibly difficult. Take the example where $\psi_{ij}(x_i, x_j) = 1[x_i x_j = 0]$. Then no two adjacent x_i, x_j can both be 1. It follows that Z counts the number of independent sets in G , which is #P complete (very hard).

Since we cannot compute Z or sample from μ exactly, we must use approximative methods, which include MCMC (Markov Chain Monte Carlo), variational inference, or diffusion models. For now, we focus on variational inference, which aims to

approximate μ by metric of distance from a family \mathcal{P} of simpler distributions (e.g. product distributions). Then the goal is the following:

$$\min_{\nu \in \mathcal{P}} KL(\nu || \mu). \quad (7)$$

Of course, if \mathcal{P} is the family of all distributions, then the minimizer is simply $\nu = \mu$ by Gibbs' inequality on the KL distance. While SOS pseudodistributions expand the space of distributions, variational inference diminishes it. The KL optimizer is hard to evaluate, but we can get around it. Notice

$$\mathbb{E}_{x \in \nu} \left[\log \frac{\nu(x)}{\mu(x)} \right] = \mathbb{E}_{x \in \nu} \left[\log \frac{\nu(x)}{\frac{1}{Z} \cdot \exp(-\mathcal{E}(x))} \right] = \mathbb{E}_{\nu}[\log \nu] + \mathbb{E}_{\nu}[\mathcal{E}(x)] - \log \frac{1}{Z} \quad (8)$$

so

$$KL(\nu || \mu) = \mathbb{E}_{\nu}[\log \nu] + \mathbb{E}_{\nu}[\mathcal{E}(x)] - \log \frac{1}{Z}. \quad (9)$$

For the purposes of minimization over ν , the last term is independent, so only the first two terms – the negative entropy and average energy respectively – are important. The sum can be written as the Gibbs free energy functional

$$G[\nu] = \mathbb{E}_{\nu}[\log \nu] + \mathbb{E}_{\nu}[\mathcal{E}(x)]. \quad (10)$$

For the Ising model,

$$G[\nu] = \mathbb{E}_{\nu} \left[\frac{\beta}{2} x^{\top} A x \right] - H(\nu). \quad (11)$$

When β is small, maximizing entropy minimizes G . When β is large, minimizing average energy minimizes G . We will see how belief propagation is a heuristic for minimizing $G[\nu]$ in the next lecture. Belief propagation is a natural dynamic programming algorithm that gives an exact answer when the graph is a tree and finds stationary points of the Lagrangian dual of $G[\nu]$ called the Bethe free energy.

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